SMT BASICS

WS 2018 / 2019 (342.208)

Armin Biere biere@jku.at
Martina Seidl martina.seidl@jku.at

Institute for Formal Models and Verification
Johannes Kepler Universität Linz

Version 2018.1

JOHANNES KEPLER UNIVERSITÄT LINZ

JOHANNES KEPLER UNIVERSITY LINZ

Satisfiability Modulo Theories (SMT)

Example

$f(x) \neq f(y) \ \wedge \ x + u = 3 \ \wedge \ v + y = 3 \ \wedge \ u = a[z] \ \wedge \ v = a[w] \ \wedge \ z = w$

- formulas in first-order logic
  usually without quantifiers, variables implicitly existentially quantified with sorted / typed symbols
  including functions / constants / predicates are interpreted
  SMT quantifier reasoning weaker than in first-order theorem proving (FO)
  much richer language compared to propositional logic (SAT)
- no need to axiomatize “theories” using axioms with quantifiers
  important theories are “built-in”: **uninterpreted functions, equality, arithmetic, arrays, bit-vectors** . . .
  focus is on decidable theories, thus fully automatic procedures
- state-of-the-art SMT solvers essentially rely on SAT solvers
  SAT solver enumerates solutions to a propositional skeleton
  propositional and theory conflicts recorded as propositional clauses
  DPLL(T), CDCL (T), read DPLL modulo theory T or CDCL modulo T
- SMT sweet spot between SAT and FO: many (industrial) applications
  standardized language SMTLIB used in applications and competitions

Buggy Program

```
int middle (int x, int y, int z) {
  int m = z;
  if (y < z) {
    if (x < y)
      m = y;
    else if (x < z)
      m = y;
  } else {
    if (x > y)
      m = y;
    else if (x > z)
      m = x;
  }
  return m;
}
```

this program is supposed to return the middle (median) of three numbers

Test Suite for Buggy Program

```
middle (1, 2, 3) = 2
middle (1, 3, 2) = 2
middle (2, 1, 3) = 1
middle (2, 3, 1) = 2
middle (3, 1, 2) = 2
middle (3, 2, 1) = 2

middle (1, 1, 1) = 1

middle (1, 1, 2) = 1
middle (1, 2, 1) = 1
middle (2, 1, 1) = 1

middle (1, 2, 2) = 2
middle (2, 1, 2) = 2
middle (2, 2, 1) = 2
```

- This black box test suite has to be generated manually.
- How to ensure that it covers all cases?
- Need to check outcome of each run individually and determine correct result.
- Difficult for large programs.
- Better use specification and check it.

Specification for Middle

let a be an array of size 3 indexed from 0 to 2

$$
\begin{aligned}
&\quad a[i] = x \wedge a[j] = y \wedge a[k] = z \\
&\wedge \\
&\quad a[0] \leq a[1] \wedge a[1] \leq a[2] \\
&\wedge \\
&\quad i \neq j \wedge i \neq k \wedge j \neq k \\
&\rightarrow \\
&\quad m = a[1]
\end{aligned}
$$

median obtained by sorting and taking middle element in the order

coming up with this specification is a manual process

Encoding of Middle Program in Logic

```
  int m = z;
  if (y < z) {
    if (x < y)
      m = y;
    else if (x < z)
      m = y;
  } else {
    if (x > y)
      m = y;
    else if (x > z)
      m = x;
  }
  return m;
}
```

$$
\begin{aligned}
& (y < z \land x < y \to m = y) \\
\land \quad & \\
& (y < z \land x \geq y \land x < z \to m = y) \\
\land \quad & \\
& (y < z \land x \geq y \land x \geq z \to m = z) \\
\land \quad & \\
& (y \geq z \land x > y \to m = y) \\
\land \quad & \\
& (y \geq z \land x \leq y \land x > z \to m = x) \\
\land \quad & \\
& (y \geq z \land x \leq y \land x \leq z \to m = z)
\end{aligned}
$$

this formula can be generated automatically by a compiler

Translating Checking of Specification as SMT Problem

let P be the encoding of the program, and S of the specification

program is correct if "$P \rightarrow S$" is valid

program has a bug if "$P \rightarrow S$" is invalid

program has a bug if negation of "$P \rightarrow S$" is satisfiable (has a model)

program has a bug if "$P \wedge \neg S$" is satisfiable (has a model)

Checking Specification as SMT Problem Example

$$
\begin{array}{ll}
(y < z \land x < y \rightarrow m = y) & \land \\
(y < z \land x \geq y \land x < z \rightarrow m = y) & \land \\
(y < z \land x \geq y \land x \geq z \rightarrow m = z) & \land \\
(y \geq z \land x > y \rightarrow m = y) & \land \\
(y \geq z \land x \leq y \land x > z \rightarrow m = x) & \land \\
(y \geq z \land x \leq y \land x \leq z \rightarrow m = z) & \land \\
a[i] = x \land a[j] = y \land a[k] = z & \land \\
a[0] \leq a[1] \land a[1] \leq a[2] & \land \\
i \neq j \land i \neq k \land j \neq k & \land \\
m \neq a[1] &
\end{array}
$$

Encoding with Linear Integer Arithmetic in SMTLIB2

```
(set-logic QF_AUFLIA)
(declare-fun x () Int) (declare-fun y () Int) (declare-fun z () Int) (declare-fun m () Int)
(assert (=> (and (< y z) (< x y) ) (= m y)))
(assert (=> (and (< y z) (>= x y) (< x z)) (= m y)))      ; fix by replacing last 'y' by 'x'
(assert (=> (and (< y z) (>= x y) (>= x z)) (= m z)))
(assert (=> (and (>= y z) (> x y) ) (= m y)))
(assert (=> (and (>= y z) (<= x y) (> x z) ) (= m x)))
(assert (=> (and (>= y z) (<= x y) (<= x z)) (= m z)))
(declare-fun i () Int) (declare-fun j () Int) (declare-fun k () Int)
(declare-fun a () (Array Int Int))
(assert (and (<= 0 i) (<= i 2) (<= 0 j) (<= j 2) (<= 0 k) (<= k 2)))
(assert (and (= (select a i) x) (= (select a j) y) (= (select a k) z)))
(assert (<= (select a 0) (select a 1) (select a 2)))
(assert (distinct i j k))
(assert (distinct m (select a 1)))
(check-sat) (get-model) (exit)
```

Checking Middle Example with Z3

```
$ z3 middle-buggy.smt2
sat
(model
  (define-fun i () Int 1)
  (define-fun a () (Array Int Int) (_ as-array k!0))
  (define-fun j () Int 0)
  (define-fun k () Int 2)
  (define-fun m () Int 2281)
  (define-fun z () Int 2283)
  (define-fun y () Int 2281)
  (define-fun x () Int 2282)
  (define-fun k!0 ((x!1 Int)) Int
    (ite (= x!1 2) 2283
    (ite (= x!1 1) 2282
    (ite (= x!1 0) 2281 2283))))
)
```

```
$ z3 middle-fixed.smt2
unsat
```

see also `http://rise4fun.com`

Encoding with Bit-Vector Logic in SMTLIB2

```
(set-logic QF_AUFBV)
(declare-fun x () (_ BitVec 32)) (declare-fun y () (_ BitVec 32))
(declare-fun z () (_ BitVec 32)) (declare-fun m () (_ BitVec 32))
(assert (=> (and (bvult y z) (bvult x y) ) (= m y)))
(assert (=> (and (bvult y z) (bvuge x y) (bvult x z)) (= m y)))    ; fix last 'y'->'x'
(assert (=> (and (bvult y z) (bvuge x y) (bvuge x z)) (= m z)))
(assert (=> (and (bvuge y z) (bvugt x y) ) (= m y)))
(assert (=> (and (bvuge y z) (bvule x y) (bvugt x z)) (= m x)))
(assert (=> (and (bvuge y z) (bvule x y) (bvule x z)) (= m z)))
(declare-fun i ()(_ BitVec 2)) (declare-fun j ()(_ BitVec 2)) (declare-fun k ()(_ BitVec 2))
(declare-fun a ()(Array (_ BitVec 2) (_ BitVec 32)))
(assert (and (bvule #b00 i) (bvule i #b10) (bvule #b00 j) (bvule j #b10)))
(assert (and (bvule #b00 k) (bvule k #b10)))
(assert (and (= (select a i) x) (= (select a j) y) (= (select a k) z)))
(assert (bvule (select a #b00) (select a #b01)))
(assert (bvule (select a #b01) (select a #b10)))
(assert (distinct i j k)) (assert (distinct m (select a #b01)))
(check-sat) (get-model) (exit)
```

Checking Middle Example with Boolector

```
$ boolector -m middle32-buggy.smt2
sat
...
2 11001101100011110101101111001001 x
3 01101101100011110101101110000001 y
4 11101011000011110101100111010001 z
5 01101101100011110101101110000001 m
28 01 i
29 00 j
30 10 k
31[00] 01101101100011110101101110000001 a
31[01] 11001101100011110101101111001001 a
31[10] 11101011000011110101100111010001 a

$ boolector middle32-fixed.smt2
unsat
```

see also `http://fmv.jku.at/boolector`

Theory of Linear Real Arithmetic (LRA)

- constants: integers, rationals, etc.
- predicates: equality $=$, disequality $\neq$, inequality $\leq$ (strict $<$) etc.
- functions: addition $+$, subtraction $-$, multiplication $\cdot$ by constant only

Example

$z \leq x - y \;\wedge\; x + 2 \cdot y \leq 5 \;\wedge\; 4 \cdot z - 2 \cdot x \geq y$

- we focus on conjunction of inequalities as in the example first
- equalities "$=$" can be replaced by two inequalities "$\leq$"
  - disequalities replaced by disjunction of strict inequalities
- combination with SAT allows arbitrary formulas (not just conjunctions)
- related to optimization problems solved in operation research (OR)
  - OR algorithms are usually variants of the classic SIMPLEX algorithm

Fourier-Motzkin Elimination Procedure by Example

$$z \leq x - y \quad \wedge \quad x + 2 \cdot y \leq 5 \quad \wedge \quad 4 \cdot z - 2 \cdot x \geq y$$

pick *pivot* variable, e.g. x, and *isolate* it on one side with coefficient 1

$$\begin{aligned}
z + y \leq x \quad &\wedge \quad x \leq 5 - 2 \cdot y \quad \wedge \quad 4 \cdot z - y \geq 2 \cdot x \\
z + y \leq x \quad &\wedge \quad x \leq 5 - 2 \cdot y \quad \wedge \quad 2 \cdot z - 0.5 \cdot y \geq x \\
z + y \leq x \quad &\wedge \quad x \leq 5 - 2 \cdot y \quad \wedge \quad x \leq 2 \cdot z - 0.5 \cdot y \qquad (1)
\end{aligned}$$

eliminate x by adding $A \leq B$ for all inequalities $A \leq x$ and $x \leq B$

$$\begin{aligned}
z + y \leq 5 - 2 \cdot y \quad &\wedge \quad z + y \leq 2 \cdot z - 0.5 \cdot y \\
z \leq 5 - 3 \cdot y \quad &\wedge \quad 1.5 \cdot y \leq z \qquad (2)
\end{aligned}$$

and same procedure with new pivot variable, e.g. z, and eliminate z

$$\begin{aligned}
1.5 \cdot y &\leq 5 - 3 \cdot y \\
y &\leq 10/9 \qquad (3)
\end{aligned}$$

(3) has (as one) solution $y = 0 \in (-\infty, 10/9]$ or $y = 1 \in (-\infty, 10/9]$

(2) then allows $z = 0 \in [0, 5]$ $z = 2 \in [1.5, 2]$

(1) then forces $x = 0$ forces $x = 3$ thus *satisfiable*

Theory of Uninterpreted Functions and Equality

- functions as in first-order (FO): sorted / typed without interpretation
- equality as single interpreted predicate
  - *congruence axiom* $\forall x, y\colon x = y \rightarrow f(x) = f(y)$
  - similar variants for functions with multiple arguments
  - always assumed in FO if equality is handled explicitly (interpreted)
- uninterpreted functions allow to abstract from concrete implementations
  - in hardware (HW) verification abstract complex circuits (e.g. multiplier)
  - in software (SW) verification abstract sub routine computation
- *congruence closure* algorithms using fast union-find data structures
  - start with all terms (and sub-terms) in different equivalence classes
  - if $t_1 = t_2$ is an asserted literal merge equivalence classes of t_1 and t_2
  - for all elements of an equivalence class check congruence axiom
    - let t_1 and t_2 be two terms in the same equivalence class
    - if there are terms $f(t_1)$ and $f(t_2)$ merge their equivalence classes
  - continue until the partition of terms in equivalence classes stabilizes
  - if asserted disequality $t_1 \neq t_2$ exists with t_1, t_2 in the same equivalence class then *unsatisfiable* otherwise *satisfiable*

Congruence Closure By Example

assume flattened structure where all sub-terms are identified by variables

$$[x \mid y \mid t \mid u \mid v]$$

$$\underbrace{x = y}_{} \wedge x = g(y) \wedge t = g(x) \wedge u = f(x,t) \wedge v = f(y,x) \wedge u \neq v$$

asserted literal $x = y$ puts x and y in to the same equivalence class

$$[x\ y \mid t \mid u \mid v]$$

$$x = y \wedge \underbrace{x = g(y) \wedge t = g(x)}_{} \wedge u = f(x,t) \wedge v = f(y,x) \wedge u \neq v$$

apply congruence axiom since x and y in same equivalence class

Congruence Closure By Example

$$[x\ y\ t \mid u \mid v]$$

$$x = y \land x = g(y) \land t = g(x) \land \underbrace{u = f(x,t) \land v = f(y,x)}_{\text{apply congruence axiom since } y,\, x \text{ and } t \text{ are all in same equivalence class}} \land u \neq v$$

$$[x\ y\ t \mid u\ v]$$

$$x = y \land x = g(y) \land t = g(x) \land u = f(x,t) \land v = f(y,x) \land u \neq v$$

u and v in the same equivalence class but $u \neq v$ asserted
thus *unsatisfiable*