

- b) ... compute the *minimum reachable DFSI (MRDFSI)* for each node.¹
- c) ... mark back edges with 'b' and cross edges with 'c'.
- d) ... mark all strongly-connected components with circles.

Comment: For the following exercises, assume that parse DAGs are implemented like presented in the lecture (e.g. see slide 121).

Exercise 35

- a) Given boolean variables x and y , find two different formulations of the binary XOR-operation $x \oplus y$ using only negation and binary conjunction.
- b) Find a DNF representation² for the parity function f over four boolean variables:
 $f(x_1, x_2, x_3, x_4) := x_1 \oplus x_2 \oplus x_3 \oplus x_4$.
- c) Given three boolean variables x_1, x_2, x_3 , draw all structurally different³ parse DAGs for representing ternary conjunction over x_1, x_2, x_3 using only binary AND-operator nodes.

Exercise 36

- a) Reformulate the binary logical operators disjunction \vee , implication \rightarrow , equivalence \leftrightarrow and XOR \oplus using only binary conjunction and negation.
- b) For each of the reformulations from a), draw a parse DAG with explicit sharing of nodes. How many nodes could be saved in each parse DAG when implementing negation as an edge attribute (i.e. encoded in the LSB of the pointer) instead of separate NOT-operator nodes? Justify your answer.
- c) Given n distinct parse DAG variable nodes v_1, v_2, \dots, v_n . Let D denote a parse DAG representing the n -ary conjunction $v_1 \wedge v_2 \wedge \dots \wedge v_n$ over all v_i . Given D , let the *level* of D be the maximal length of a path from the root of D to any v_i in D . What is the *maximal* and *minimal* level of any possible D ? Justify your answer.

¹Specify this value *before* it is reset to INF by `mrdfs_i[M] = INF` near the bottom of slide 110.

²Note that this exercise can be solved without constructing the truth table of f .

³Two parse DAGs are structurally different if the pointers to children in `child[0]` or `child[1]` are different.